

CRYPTOGRAPHY SECURITY TO INTERNET OF THING

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ABSTRACT-Increase in the use of the network there is a need for the data to be secured, to make the data secure there exist number of algorithms like AES (Advance Encryption Standard) and IDEA (International Data Encryption Algorithm). It secures data using cryptography. Internet of thing (IOT) to the cryptography is new domain which is emerging widely. IOT is defined as the controlling any data from anywhere using internet. Thus, IOT require security. In this paper discuss about how data can be secured for IOT using the methods of cryptography[1].

KEYWORDS-Advance Encryption Standard, Internet of thing, Data encryption, cryptography.

INTRODUCTION-IOT security[2][3] involves securing the data and securing the data and uploading it to cloud. Encrypting the data's using AES-128 algorithm[4]. This involves following steps-

S-box, Shift rows, Mix Column, Add round key.

On following the steps for 10 rounds we get the encrypted data. This encrypted data is uploaded to the cloud using IOT development board. Here Intel Galileo board is used. Similarly to retrieve the data inverse AES-128 algorithm is used which involves following steps-

Inverse S-box, Inverse Shift rows, Inverse Mix Column, Inverse Add round key.

On following the steps for 10 rounds we get the decrypted data i.e. original data. Fig.1 gives the block diagram overview.

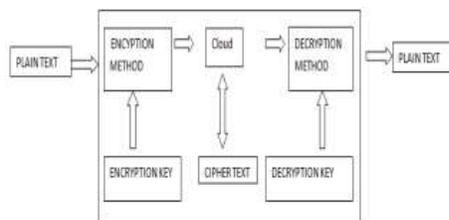


Fig. 1. Block Diagram of IoT Security

AES-128 ALGORITHM- Advanced Encryption Algorithm is also known as Rijndael was established by U.S National Institute of Standards and Technology (NIST) in 2001[5]. AES uses 3 types of data which are- 128 bits with 10 rounds, 192 bits with 12 rounds, 256 bits with 14 rounds. AES-128 uses 10 rounds each with 4 steps[6]. Thus, data form a 4x4 matrix with each

elements being 8bits. Data input is a plain text giving an encrypted data while for the decryption input data is a cipher data with original data as output.

S- BOX- In this each element in the plain text is replaced by the element of the s-box, where s-box is a 16x16 matrix with each element with 8 bits. Replacement of the data is done in such a way that first8 bits represent row and next 8 bit represent column.

63	7c	77	7b	f2	6b	6f	c5	30	01	67	2b	fe	d7	ab	7d
ca	82	c9	7d	fa	59	47	f0	ad	d4	a2	af	9c	a4	72	c0
b7	fd	93	26	36	3f	f7	cc	34	a5	e5	f1	71	d8	31	15
04	c7	23	c3	18	96	05	9a	07	12	80	e2	ab	27	b2	75
09	83	2c	1a	1b	6e	5a	a0	52	3b	d6	b3	29	e3	2f	84
53	d1	00	ed	20	fc	b1	5b	6a	cb	be	39	4a	4c	58	cf
d0	ef	aa	fb	43	4d	33	85	45	f9	02	7f	50	3c	9f	a8
51	a3	40	8f	92	9d	38	f5	bc	b6	da	21	10	ff	f3	d2
cd	0c	13	ec	5f	97	44	17	c4	a7	7e	3d	64	5d	19	73
60	91	4f	dc	22	2a	90	88	46	ee	b8	14	de	5e	0b	db
e0	32	3a	0a	49	06	24	5c	e2	d3	ac	62	91	95	e4	79
e7	c8	37	6d	8d	d5	4e	a9	6c	56	f4	ea	65	7a	ae	08
ba	78	25	2e	1c	a6	b4	c6	e8	dd	74	1f	4b	bd	8b	8a
70	3e	b5	66	48	03	26	0a	61	35	57	b9	86	e1	1d	9e
e1	f8	98	11	69	d9	8e	94	9b	1e	87	e9	ce	55	28	df
8c	a1	89	0d	bf	e6	42	68	41	99	2d	0f	b0	54	bb	16

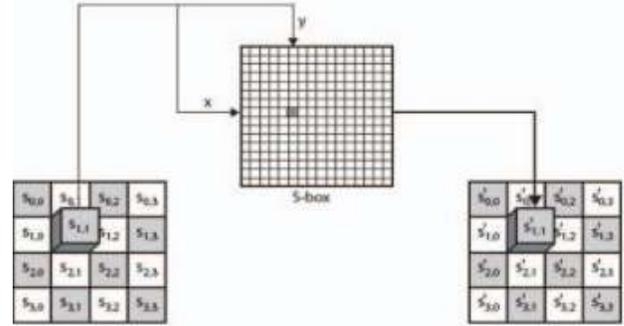


Fig. 2. S-Box

- SHIFT ROWS- This step is applied to the row of the matrix, in this first row remain unchanged while second row get left shift by 1 cyclically, third row get left shift by 2 cyclically, and so on. Which give a 4x4 shifted matrix with each elements of 8 bits.

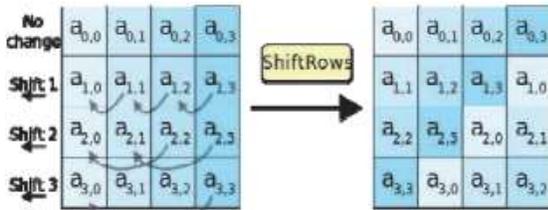


Fig. 4. Shift Row Overview

- MIX COLUMN- This step involves a matrix multiplication with a polynomial. Which give a 4x4 matrix of with each element of 8 bits. For 0x01 multiplication there is no change to the value while multiplication with 0x02 shift data to left by 1 bit, if original data has high bit at MSB (Most significant bit) then sifted data need to be XORed with 0x1b and so on

$$c(x) = \begin{bmatrix} 0x02 & 0x03 & 0x01 & 0x01 \\ 0x01 & 0x02 & 0x03 & 0x01 \\ 0x01 & 0x01 & 0x02 & 0x03 \\ 0x03 & 0x01 & 0x01 & 0x02 \end{bmatrix}$$

Fig. 5. c(x) matrix

$$\begin{bmatrix} 02 & 03 & 01 & 01 \\ 01 & 02 & 03 & 01 \\ 01 & 01 & 02 & 03 \\ 03 & 01 & 01 & 02 \end{bmatrix} \begin{bmatrix} s_{0,0} & s_{0,1} & s_{0,2} & s_{0,3} \\ s_{1,0} & s_{1,1} & s_{1,2} & s_{1,3} \\ s_{2,0} & s_{2,1} & s_{2,2} & s_{2,3} \\ s_{3,0} & s_{3,1} & s_{3,2} & s_{3,3} \end{bmatrix} = \begin{bmatrix} s'_{0,0} & s'_{0,1} & s'_{0,2} & s'_{0,3} \\ s'_{1,0} & s'_{1,1} & s'_{1,2} & s'_{1,3} \\ s'_{2,0} & s'_{2,1} & s'_{2,2} & s'_{2,3} \\ s'_{3,0} & s'_{3,1} & s'_{3,2} & s'_{3,3} \end{bmatrix}$$

Fig. 6. Mix Column Step

- ADD ROUND KEYS- This generate key having XOR operations of the mix column matrix with generated for corresponding round gives the output for the current rounds.

The output of this round steps after 10 rounds provides the encrypted data. The key for each round is derived from the original key using Rijndael’s key schedule.

INVERSE AES-128 ALGORITHM- Input to the inverse algorithm is the encrypted data the output is plain text or original data. Inverse AES-128 algorithm involves the following steps-

- INVERSE S-BOX- it works same as s-box, steps provide and output of 4x4 matrix with each element of 8 bits.

52	09	6a	d5	30	36	a5	38	bf	40	a3	9e	81	f3	d7	fb
7c	e3	39	82	9b	2f	ff	87	34	8e	43	44	c4	de	e9	cb
54	7b	94	32	a6	c2	23	3d	ee	4c	95	0b	42	fa	c3	4e
08	2e	a1	66	28	d9	24	b2	76	5b	a2	49	6d	8b	d1	25
72	f8	f6	64	8f	68	98	16	d4	a4	5c	cc	5d	65	b6	92
6c	70	48	50	fd	ed	b9	da	5e	15	46	57	a7	8d	9d	84
90	d8	ab	00	8c	bc	d3	0a	f7	e4	58	05	b8	b3	45	06
d0	2c	1e	8f	ca	3f	0f	02	c1	af	bd	03	01	13	8a	6b
3a	91	11	41	4f	67	dc	ea	97	f2	cf	ce	f0	b4	e6	73
56	ac	74	22	e7	ad	35	85	e2	f9	37	e8	1c	75	df	6e
47	f1	1a	71	1d	29	c5	89	6f	b7	62	0e	aa	18	be	1b
fc	56	3e	4b	c6	d2	79	20	9a	db	c0	fe	78	cd	5a	f4
1f	dd	a8	33	88	07	c7	31	b1	12	10	59	27	80	ec	5f
60	51	7f	a9	19	b5	4a	0d	2d	e5	7a	9f	93	c9	9c	ef
a0	e0	3b	4d	ae	2a	f5	b0	c8	eb	bb	3c	83	53	99	61
17	2b	04	7e	ba	77	d6	26	e1	69	14	63	55	21	0c	7d

Fig. 7. Inverse S-box

- INVERSE SHIFT ROW- In this first row is kept same, while second rows is shifted towards 1 by right cyclically. The third row is shifted towards right by 2 cyclically and so on.
- INVERSE MIX COLUMN- it is similar to the Mix column step as shown in figure.
- INVERSE ADD ROUND KEY- In this a XOR of keys for the corresponding round with the inverse mix column matrix gives an output. Again key generation involves Rijndael key schedule.

$$c(x) = \begin{bmatrix} 0x0e & 0x0b & 0x0d & 0x09 \\ 0x09 & 0x0e & 0x0b & 0x0d \\ 0x0d & 0x09 & 0x0e & 0x0b \\ 0x0b & 0x0d & 0x09 & 0x0e \end{bmatrix}$$

Fig. 8. Inverse Mix Column Step

APPLICATION- Bird Hit Probability-In this using a camera we can capture the image of bird. Using image processing we can detect the position of bird using x, y and z coordinates. Thus this plain data is encrypted and uploaded to cloud. At other end, the received data is decrypted and use those data to avoid the collision of airplane and birds.

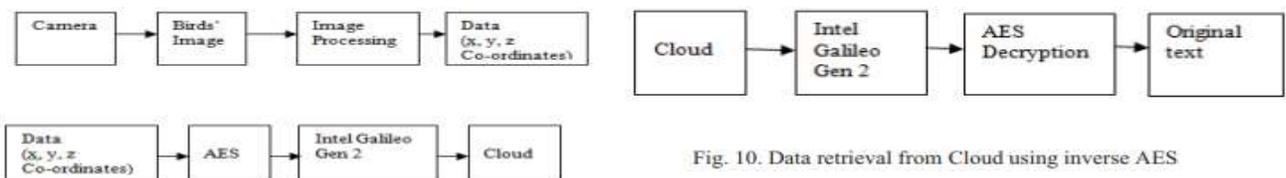


Fig. 10. Data retrieval from Cloud using inverse AES

Fig. 9. Uploading encrypted data to Cloud

RESULTS AND APPLICATION-

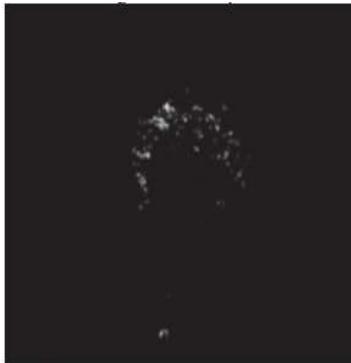


Fig. 11. Image when bird is detected

position x	y	z
433.000000	183.000000	42.069483
483.000000	483.000000	39.698060
423.000000	297.000000	68.883957
289.000000	277.000000	164.477266
485.000000	217.000000	73.164200
431.000000	189.000000	52.006460
427.000000	195.000000	56.302753
521.000000	367.000000	15.264330
517.000000	367.000000	13.416408
421.000000	187.000000	50.477710
421.000000	139.000000	15.132746
517.000000	385.000000	12.030404
435.000000	291.000000	145.165421
395.000000	217.000000	67.683003
399.000000	203.000000	60.000000
447.000000	185.000000	41.400402
515.000000	369.000000	13.152946
449.000000	183.000000	37.121422

Fig. 12. Data in terms of x, y and z co-ordinates

CONCLUSION AND FUTURE SCOPE- The data encryption and decryption is done for 128 bits. To make the data more secure. In bird hit probability, since the movement of the bird is not stationary, capturing the image, process it, encrypt it and upload it to cloud would consume time. To overcome this black box can be made which receive data as a plain text and give the encrypted output which can be used for many TOI applications.

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